

# EEL 4713 Assignment #1

Section 1: Done

Section 2: ch1 -

- |                           |                          |
|---------------------------|--------------------------|
| 1.1.1) Servers            | 2) Petabyte              |
| 3.) Super computer        | 4.) virtual worlds       |
| 5.) RAM (12)              | 6.) CPU (13)             |
| 7.) data centers          | 8.) multicore processors |
| 9.) low-end servers       | 10.) embedded computer   |
| 11.) VHDL                 | 12.) desktop computers   |
| 13.) compiler             | 14.) assembler           |
| 15.) Cobol                | 16.) machine language    |
| 17.) instruction          | 18.) Fortran             |
| 19.) assembly language    | 20.) operating system    |
| 21.) application software | 22.) bit                 |
| 23.) systems software     | 24.) C                   |
| 25.) high-level language  | 26.) Terabyte            |

1.7 →

1.) Ratios	clock	1.28	1.5625	2.64	3.03	10	1.8	0.74
	power	1.2424	1.195	2.06	2.88	2.587	1.3678	0.9223

Geometric means:

clock → 2.15107
power → 1.61573

- 2) largest clock change ⇒ 10 (Pentium Pro → Pentium 4 Willamette)  
 largest power change ⇒ 2.88 (Pentium to Pentium Pro)

3) clock ratio  $\rightarrow \frac{2667}{12.9} = 213.36$

power ratio  $\rightarrow \frac{95}{3.3} = 28.7879$

4) Power = (Capacitive load)  $\cdot$  (Voltage)<sup>2</sup>  $\cdot$  (Frequency)  $C = \frac{P}{V^2 \cdot F}$

	Capacitive load
80286	10.56 nF
80386	10.25 nF
80486	7.84 nF
Pentium	6.1212 nF
Pentium Pro	13.36 nF
Pentium 4 Willamette	12.2939 nF
Pentium 4 Prescott	18.311 nF
Core 2 Kentsfield	29.4385 nF

5) Voltage Ratios

$5 \rightarrow 3.3 = 0.66$

$3.3 \rightarrow 1.75 = 0.53$

$1.75 \rightarrow 1.25 = 0.7143$

$1.25 \rightarrow 1.1 = 0.88$

largest change = 0.53

Pentium Pro to Pentium 4 W

6) geometric mean = 0.684866

$\$t_1 =$  some kind of count value

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## Section 3: Lab 1.s

a) let  $a_0 = x = 6$   
 $a_1 = y = 4$   
 $a_2 = z = 8$

does  
 $\$t_0 = 6 \cdot 4$

```
main: addi $a0, $0, 6
      addi $a1, $0, 4
      addi $a2, $0, 8
      addi $t0, $0, 0
      addi $t1, $a1, 0 ( $t_1 = a_1$ )
```

```
loop1: add $t0, $t0, $a0 ( $t_0 = t_0 + a_0$ )
        addi $t1, $t1, -1
        bne $t1, $0, loop1
```

32  $\$t_2 = 8 + (6 \cdot 4) \rightarrow$   
16  $\$t_3 = (6 \cdot 4) - 8 \rightarrow$   
 $\$v_0 = 0$

```
add $t2, $t0, $a2
sub $t3, $t0, $a2
add $v0, $0, $0
```

performs  $\frac{((6 \cdot 4) - 8)}{2}$   
 $\$v_0 = (8 + (6 \cdot 4)) \cdot \frac{1}{2}$   
 $= 256$

```
loop2: add $v0, $v0, $t2
        addi $t3, $t3, -2
        bne $t3, $0, loop2
```

so let  $\$a_0 = x, \$a_1 = y, \$a_2 = z, \$v_0 = r$

this code performs the operation:

$$r = \frac{[z + (x \cdot y)] + [(x \cdot y) - z]}{2}$$

for  $x=6, y=4$   
for  $z=8$   
 $r=256$

b) which registers are used to hold operands & results?

I'm assuming for example `add $t2, $t0, $a2`  
↑ arguments  
↓ result  
operands (also immediate values)

i) for the overall process `$a0`, `$a1`, & `$a2` hold the arguments to the function and `$v0` stores the result

therefore, overall

operands: <code>\$a0</code> , <code>\$a1</code> , <code>\$a2</code> , <code>-1</code> , <code>-2</code>
results: <code>\$v0</code>

ii) however, some registers hold intermediate values that are both results and operands for later instructions

- `$t0` holds the result for a multiply but is used as an argument for addition & subtraction instructions later
- `$t1` & `$t3` act as count values for the multiply operations and are used as results & operands
- `$t2` stores the result of an addition, but is also the operand in a multiply later on

when including intermediate steps:

operands: <code>\$a0</code> , <code>\$a1</code> , <code>\$a2</code> , <code>-1</code> , <code>-2</code>
results: <code>\$v0</code>
both: <code>\$t0</code> , <code>\$t1</code> , <code>\$t2</code> , <code>\$t3</code>